

## 11.5 *Minimum Spanning Trees*

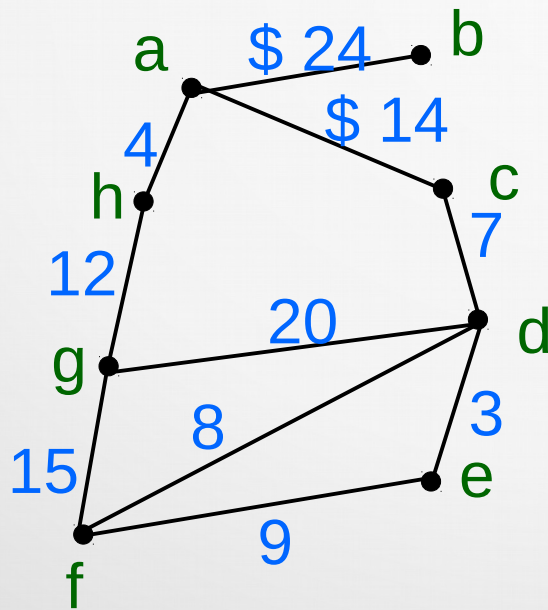
Let's consider graphs with weighted edges.

**[Def]** A *minimum spanning tree* in a connected weighted graph is a spanning tree that has the smallest possible sum of weights of its edges.

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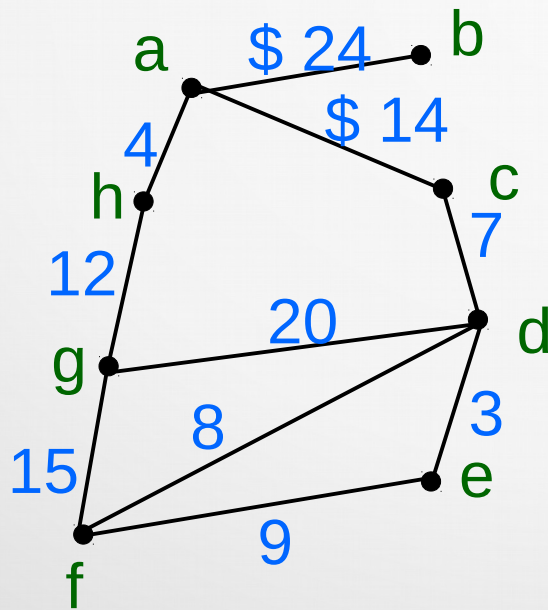


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$$\$15 + \$12 + \$4 = \$31$$

$$\$9 + \$3 + \$7 + \$14 = \$33$$

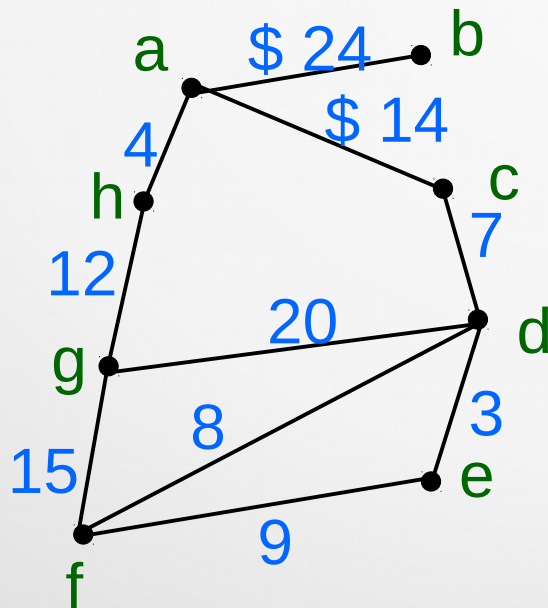
$$\$8 + \$7 + \$14 = \$29$$

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$$f \rightarrow d \rightarrow c \rightarrow a$$

...



## 11.5 *Minimum Spanning Trees*

### Prim's algorithm

**procedure** *Prim*( $G$ : weighted connected undirected graph with  $n$  vertices)

$T :=$  a minimum weight edge

**for**  $i := 1$  to  $n-2$

$e :=$  an edge of minimum weight incident to a vertex in  $T$  and not forming a simple circuit in  $T$  if added to  $T$

$T := T$  with added edge  $e$

**return**  $T$

We will assume that the edges are ordered when we need to choose between two or more edges with the same weights.

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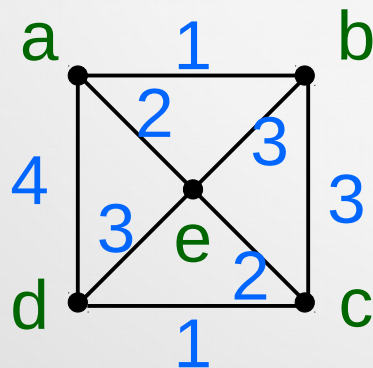
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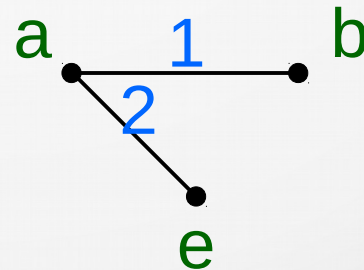
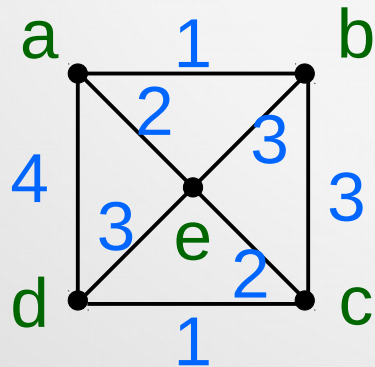
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**procedure** *Prim*( $G$ : weighted conn. undir.graph with  $n$  vertices)  
T := a minimum weight edge  
**for**  $i := 1$  to  $n-2$      $i = 1$   
    →  $e :=$  an edge of minimum weight incident to a vertex  
    in T and not forming a simple circuit in T if added to T  
    T := T with added edge  $e$   
**return** T

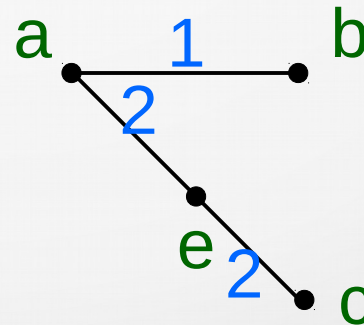
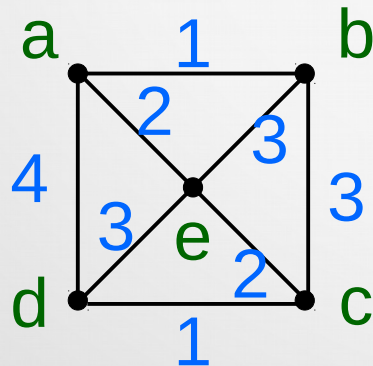




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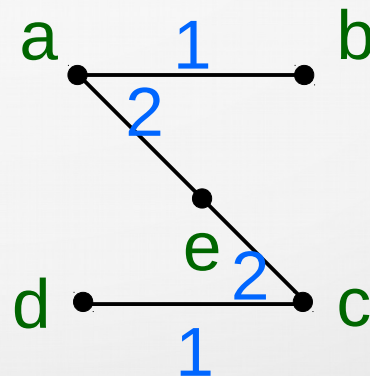
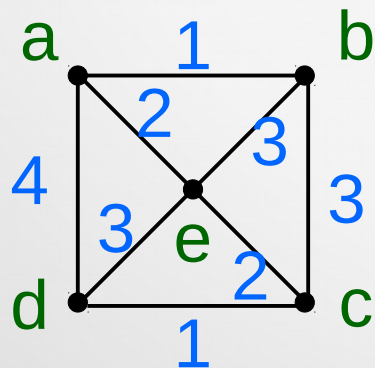
**procedure** *Prim*( $G$ : weighted conn. undir.graph with  $n$  vertices)  
   $T :=$  a minimum weight edge  
  **for**  $i := 1$  to  $n-2$      $i = 2$   
     $\rightarrow e :=$  an edge of minimum weight incident to a vertex  
      in  $T$  and not forming a simple circuit in  $T$  if added to  $T$   
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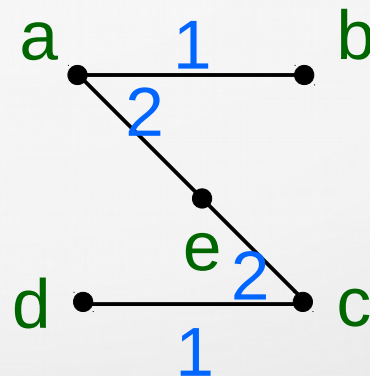
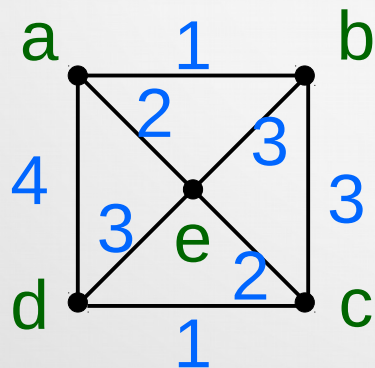
$T :=$  a minimum weight edge

**for**  $i := 1$  to  $n-2$      $i = 3$

$e :=$  an edge of minimum weight incident to a vertex  
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$T := T$  with added edge  $e$

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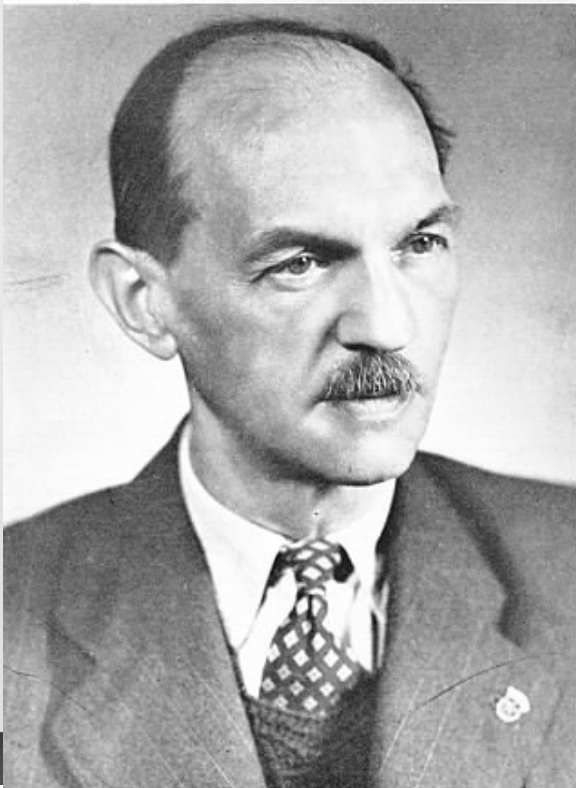


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## Prim's algorithm

Historical information:

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During his career at Bell Laboratories, Robert Prim along with coworker Joseph Kruskal developed two different algorithms for finding a minimum spanning tree in a weighted graph



# 11.5 Minimum Spanning Trees

## Kruskal's algorithm

The second algorithm up for discussion was discovered by Joseph Kruskal in 1956, although the basic ideas it uses were described much earlier.



**procedure** *Kruskal*( $G$ : weighted connected undirected graph with  $n$  vertices)

$T :=$  empty graph

**for**  $i := 1$  to  $n-1$

$e :=$  any edge in  $G$  with smallest weight that does not form a simple circuit when added to  $T$

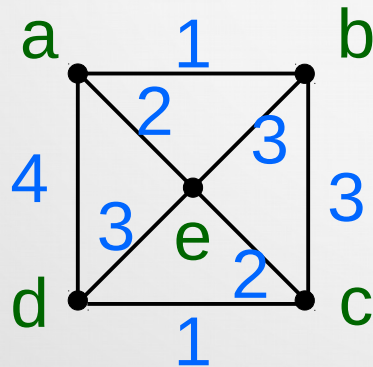
$T := T$  with added  $e$

**return**  $T$

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## Kruskal's algorithm

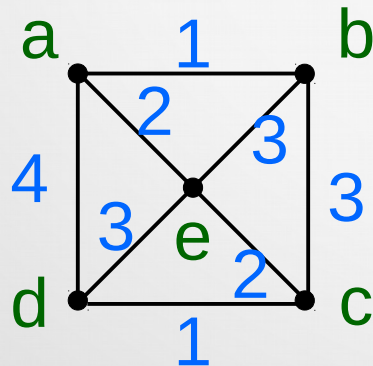
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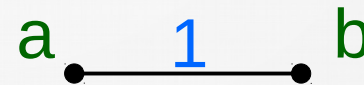
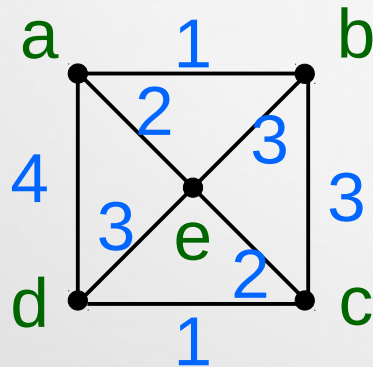




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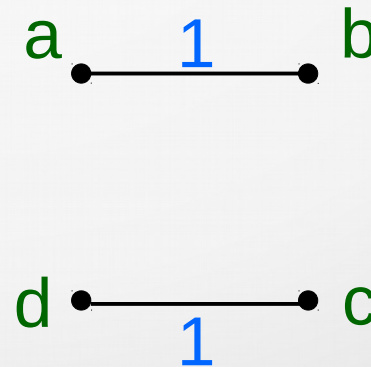
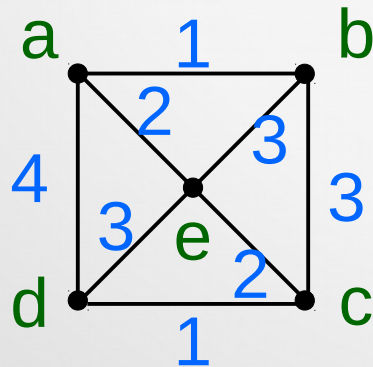
```
procedure Kruskal( $G$ : weighted conn. undir.graph with  $n$  vertices)  
   $T :=$  empty graph  
  for  $i := 1$  to  $n-1$      $i := 1$   
     $\rightarrow e :=$  any edge in  $G$  with smallest weight that does not form a  
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       $T := T$  with added  $e$   
  return  $T$ 
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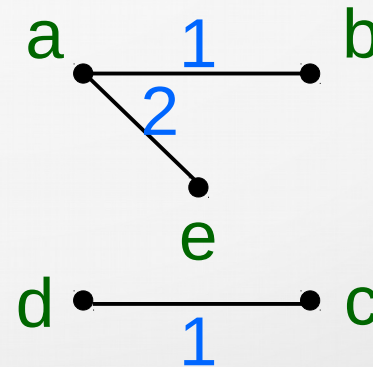
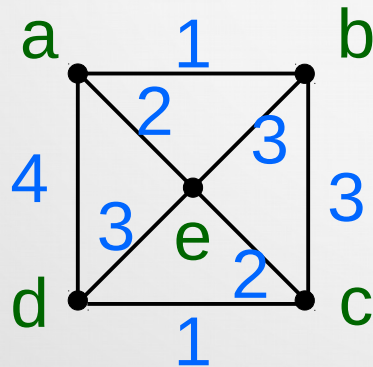
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   $T :=$  empty graph  
  for  $i := 1$  to  $n-1$        $i := 2$   
     $\rightarrow e :=$  any edge in  $G$  with smallest weight that does not form a  
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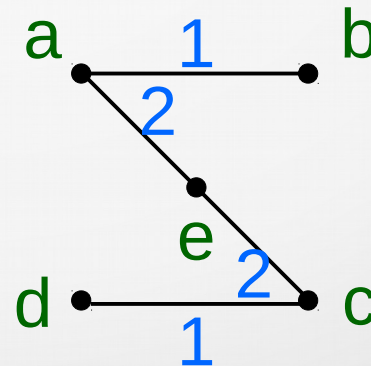
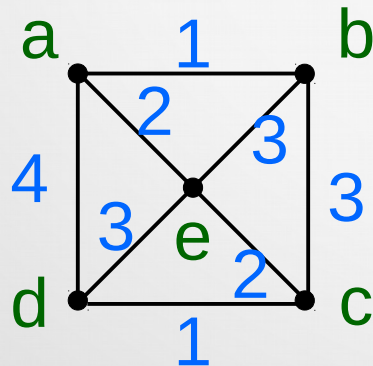
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   $T :=$  empty graph  
  for  $i := 1$  to  $n-1$        $i := 3$   
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   $T :=$  empty graph  
  for  $i := 1$  to  $n-1$        $i := 4$   
     $\rightarrow e :=$  any edge in  $G$  with smallest weight that does not form a  
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     $T := T$  with added  $e$   
→ return  $T$ 
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