

OUTLINE

- 1 CHAPTER 14: GRAPHS
 - Graph Data Structures

GRAPHS

Graphs can represent airlines, electrical circuits, or computer networks.

A Graph G will consist of:

- A set V of **vertices** (nodes, points).
(*Cities, circuit connections, computers*).
We will use V to mean the **number** of vertices as well (mathematicians use *cardinality* notation, $|V|$).
- A set E of **edges** (lines connecting vertices).
(*Air lanes, elements in a circuit, computer connections in a network*).
We will use E to mean the **number** of edges as well.

GRAPHS

- A **path** is a series of edges connecting two vertices.
- In an **undirected graph** edges are “two-way streets”
- A **connected graph** is one in which every pair of vertices is connected by a path.
- A **complete graph** is one in which every pair of vertices is connected by an edge.
- Two vertices are **adjacent** if there is an edge connecting them.
- A **cycle** in a directed graph is a loop formed by adjacent vertices.

GRAPHS

In directed graphs:

- edges are “one-way streets” beginning at one vertex and ending at another.
- **in-degree** of a vertex = # of edges ending at that vertex.
- **out-degree** of a vertex = # of edges beginning at that vertex.
- A **directed acyclic graph (DAG)** is a directed graph containing no cycles.
- Vertex **B** is **adjacent** to vertex **A** if there is an edge from **A** to **B**.
- *Example:* A **tree** is a special type of DAG.

GRAPHS

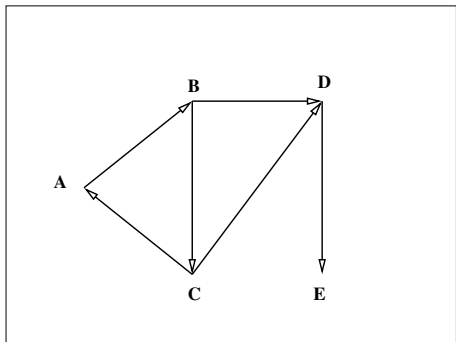
A directed graph example

in-degree of **A** =

out-degree of **C** =

in-degree of **D** =

Is there a cycle?



GRAPHS

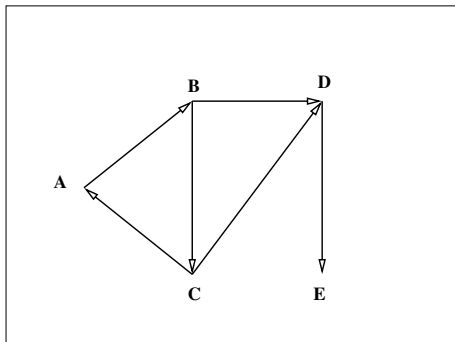
A directed graph example

in-degree of **A** = 1

out-degree of **C** = 2

in-degree of **D** = 2

Is there a cycle? Yes,
 $A \rightarrow B \rightarrow C \rightarrow A$ is a
cycle



GRAPHS

- A graph is **dense** if it has many edges connecting vertices.
- A graph is **sparse** if it has much less than the maximum possible number of edges.
- The best implementation of a graph depends on how sparse it is.
- Two commonly used data structures to represent graphs are *adjacency matrix* and *adjacency list*.

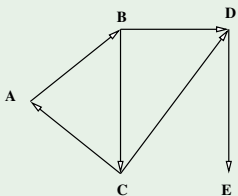
REPRESENTING GRAPHS

ADJACENCY MATRICES

- An adjacency matrix has rows and columns of zeros and ones. 1 in column i , row j means that an edge connects vertex i with vertex j (i.e. vertices i and j are adjacent).
- An adjacency matrix is used to implement a dense graph.
- It requires $\Theta(v^2)$ time to find all the edges (by checking every entry in the matrix).

REPRESENTING GRAPHS

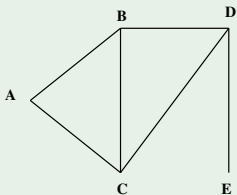
ADJACENCY MATRICES (DIRECTED GRAPH)



	A	B	C	D	E
A	0	1	0	0	0
B	0	0	1	1	0
C	1	0	0	1	0
D	0	0	0	0	1
E	0	0	0	0	0

REPRESENTING GRAPHS

ADJACENCY MATRICES (UNDIRECTED GRAPH)

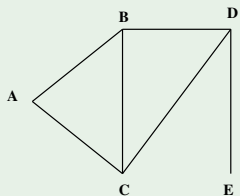


	A	B	C	D	E
A	0	1	1	0	0
B	1	0	1	1	0
C	1	1	0	1	0
D	0	1	1	0	1
E	0	0	0	1	0

The **matrix is symmetric** (entry at **row i , column j** is the same as at **row j , column i**), hence we need only half of the matrix to represent a graph (using diagonal to split it).

REPRESENTING GRAPHS

ADJACENCY MATRICES (UNDIRECTED GRAPH)

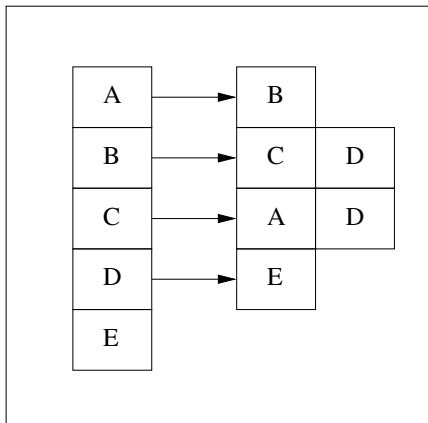
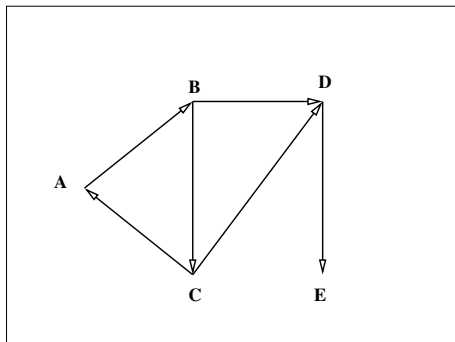


	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>
<i>A</i>	1	1	0	0
<i>B</i>		1	1	0
<i>C</i>			1	0
<i>D</i>				1

ADJACENCY LISTS

- An **adjacency list** gives each vertex an attribute which is a list of all the vertices adjacent to it.
- To represent a sparse graph, an **adjacency list** is more economical, since it only indicates where the edges are, not where they aren't.
- An adjacency list uses time $\Theta(V * E)$ to find all edges.

ADJACENCY LISTS



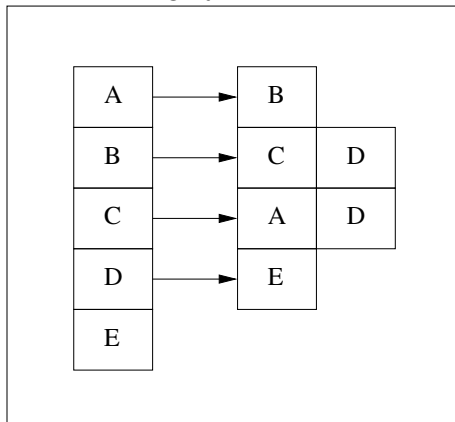
ADJACENCY LISTS

Implementation of adjacency lists in Python:

- A list of lists.
- A dictionary.

ADJACENCY LISTS: USING PYTHON LIST

Let's assume that the graph is weighted, and the weight of each edge is 1, then using Python list we can have the following:

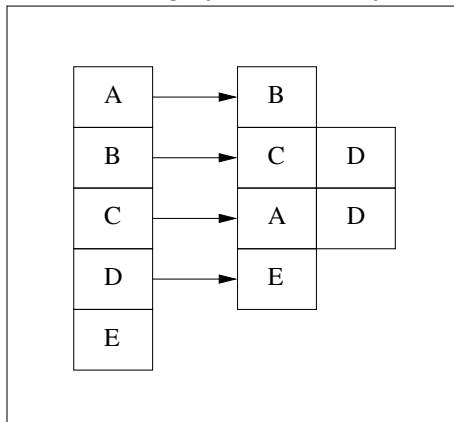


```

g = [
    ['A', [( 'B', 1)]],
    ['B', [( 'C', 1), ( 'D', 1)]],
    ['C', [( 'A', 1), ( 'D', 1)]],
    ['D', [( 'E', 1)]],
    ['E', []]]
  
```

ADJACENCY LISTS: USING PYTHON DICTIONARY

Let's assume that the graph is weighted, and the weight of each edge is 1, then using Python dictionary:



```
g = {
    'A': {'B':1},
    'B': {'C':1, 'D':1},
    'C': {'A':1, 'D':1},
    'D': {'E':1},
    'E': {}}
```


ADJACENCY LISTS: C++

Implementation of adjacency lists in C++:

- For static graphs (do not change): a two-dimensional array.
- For dynamic graphs: a list of lists (linked-list implementation).

ADJACENCY MATRIX VS ADJACENCY LIST

ADJACENCY MATRIX VS ADJACENCY LIST

- graph is dense \rightarrow the adjacency matrix representation is preferred.
- graph is sparse \rightarrow the adjacency list representation is preferred.
- Most graphs in real-world applications are sparse, hence the adjacency list representation is more commonly used.
- Using matrix representation to find all edges from a vertex we will need to examine V entries,
- Using list representation to find all edges from a vertex we will need to examine only the actual edges originating from the vertex.

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