

Short Answer

8. a)

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n = input('enter n: ') 1 operation
for i in range(n):     n iterations
    x = 2 * n           2 operations (math *, assignment)
    while x > 1:       log2 x + 1 iterations (because x is reduced
                        at each iteration - it becomes half of the
                        previous x's value - like in binary search)
        x = x/2

```

$$T(n) = 1 + n(2 + \log n + 1) = n \log n + 3n + 1$$

$T(n) = \Theta(n \log n)$ - asymptotic running time

b) $T(n) = \Theta(n)$

```

n = input('enter n: ') 1 operation
total = 0               1 operation
for i in range(n):     n iterations
    for j in range(10000): 10 000 iterations
        total += j        1 operation
print total            1 operation

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$$T(n) = 3 + 10000n$$

Since the asymptotic running time is discussed when $n \rightarrow \infty$, 10000 becomes insignificantly small there, hence $T(n) = \Theta(n)$

c) $T(n) = \Theta(n^2)$

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total = 0               1 operation
n = input('enter n: ') 1 operation
for i in range(2*n):
    for j in range(i,n): 2n + (2n-1) + (2n-2) + (2n-3) + ... + 2 + 1
                        iterations
        total += j      1 operation
for j in range(n):     n iterations
    total += j         1 operation

```

$$T(n) = 2 + 2n^2 + n + n = 2n^2 + 2n + 1, \text{ hence } T(n) = \Theta(n^2)$$

$$1 + 2 + 3 + \dots + 2n = \frac{2n(1+2n)}{2} = 2n^2 + n$$