

# OUTLINE

- 1 CHAPTER 6: RECURSION
  - Recursive Definitions
  - Simple Recursive Examples
  - In-Class Work

# RECURSIVE DEFINITIONS

## A FUNCTION CAN CALL ITSELF

- A **recursive** definition of a function is one which makes a function call to the function being defined.
- The function call is then a **recursive** function call.
- A definition is **circular** if it leads to an infinite sequence of function calls.
- To prevent this, *the function must call itself with a parameter smaller than the one it is using.*

The function must test for when the parameter has reached the minimum size (the **base** case): this must be handled without a recursive call.

# RECURSIVE DEFINITIONS

## THE CALL STACK

The function call stack can handle recursive functions easily. There is no reason why a function can't push an activation record onto the call stack with variables for the current function while calling that same function. The earlier version of that function will resume when the recursive call is completed. When the base case is finally met, there will be no further recursive calls, and no further activation records will be pushed onto the stack.

Without a base case, the stack would overflow, producing a run-time error.

## RECURSIVE DEFINITIONS

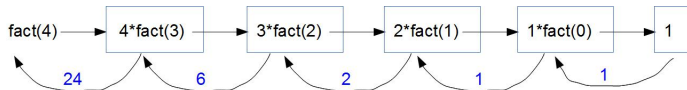
## THE FACTORIAL FUNCTION

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ n(n-1)! & \text{otherwise} \end{cases}$$

## THE FACTORIAL FUNCTION - USING PYTHON

```
def fact(n):  
    if n == 0:  
        return 1  
    else:  
        return n * fact(n - 1)
```

## RECURSIVE DEFINITIONS



## THE FACTORIAL FUNCTION - USING PYTHON

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def fact(n):  
    if n == 0:  
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```

See [fact.py](#)

# STRING REVERSAL

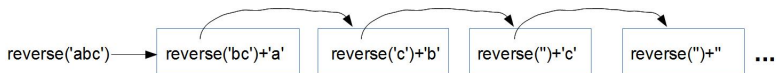
## CIRCULAR DEFINITION

```
def reverse(s):  
    return reverse(s[1:]) + s[0]
```

## STRING REVERSAL

## CIRCULAR DEFINITION

```
def reverse(s):  
    return reverse(s[1:]) + s[0]
```



# STRING REVERSAL

## DEFINITION WITH BASE CASE

```
def reverse(s):  
    if s == "":  
        return s  
    else:  
        return reverse(s[1:]) + s[0]
```

See [reverse.py](#)



# ANAGRAMS

## ANAGRAMS

An **anagram** of a word is another word spelled using the same letters but rearranged. Rearrangements are also called **permutations**. For example: **TORSO** is an anagram for **ROOST**. A recursive strategy to produce all anagrams of a given word is:

- remove the first letter from the word.
- for all anagrams of the smaller word, insert it in all possible positions.

How many permutations of the letters in TORSO is there?

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- remove the first letter from the word.
- for all anagrams of the smaller word, insert it in all possible positions.

How many permutations of the letters in TORSO is there?  
 $5! = 120$

## ANAGRAMS

## ANAGRAMS USING RECURSION

```
def anagrams(s):  
    if s == "":  
        return [s]  
    else:  
        ans = []  
        for w in anagrams(s[1:]):  
            for pos in range(len(w)+1):  
                ans.append(w[:pos]+s[0]+w[pos:1])  
        return ans
```

See [anagrams.py](#)

# FAST EXPONENTIATION

NAIVE ITERATION IS  $\Theta(n)$

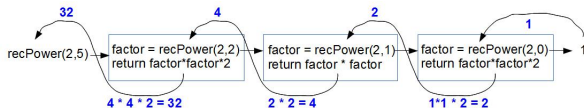
```
# power.py
def loopPower(a, n):
    ans = 1
    for i in range(n):
        ans = ans * a
    return ans
```

## FAST EXPONENTIATION

DIVIDE AND CONQUER RECURSION IS  $\Theta(\log n)$ 

```
# power.py
def recPower(a, n):
    if n == 0:
        return 1
    else:
        factor = recPower(a, n // 2)
        if n % 2 == 0:
            return factor * factor
        else:
            return factor * factor * a
```

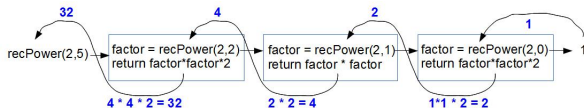
## FAST EXPONENTIATION



```
# power.py
```

```
def recPower(a, n):
    if n == 0:
        return 1
    else:
        factor = recPower(a, n // 2)
        if n % 2 == 0:
            return factor * factor
        else:
            return factor * factor * a
```

## FAST EXPONENTIATION



Let's compare the number of multiplications performed while using recursive definition of power function and naive definition of power function:

naive (running time  $\Theta(n)$ ): 4 multiplications

recursive (running time  $\Theta(\log n)$ ): 5 multiplications

If we try  $2^{10}$ , then

naive: 9 multiplications

recursive: 6 multiplications

## BINARY SEARCH

## ITERATION

```
def search(items, target):  
    low = 0  
    high = len(items) - 1  
    while low <= high:  
        mid = (low + high) // 2  
        item = items[mid]  
        if target == item:  
            return mid  
        elif target < item:  
            high = mid - 1  
        else:  
            low = mid + 1  
    return -1
```



## BINARY SEARCH

## PSEUDOCODE USING RECURSION

Algorithm: binary search

```
-- search for x in nums[low]...nums[high]
```

```
  if low > high
```

```
    x is not in nums
```

```
  mid = (low + high) // 2
```

```
  if x == nums[mid]:
```

```
    x is at mid position
```

```
  elif x < nums[mid]
```

```
    binary search for x in nums[low]...nums[mid-1]
```

```
  else
```

```
    binary search for x in nums[mid+1]...nums[high]
```

## BINARY SEARCH

## PYTHON CODE USING RECURSION

```
def search(items, target):
    return recBinSearch(target, items, 0, len(items)-1)
def recBinSearch(x, nums, low, high):
    if low > high:
        return -1
    mid = (low + high) // 2
    item = nums[mid]
    if x == item:
        return mid
    elif x < item:
        return recBinSearch(x, nums, low, mid-1)
    else:
        return recBinSearch(x, nums, mid+1, high)
```

## IN-CLASS WORK

- 1 Show pictorial representation of the call `recPower(3,7)`.
- 2 Figure out exactly how many multiplications does `recPower(3,7)` do.
- 3 Write and test a recursive function `Maximum` to find the largest number in a list.  
*Hint:* the maximum is the larger of the first item and maximum of all the other items.